ANALYSIS OF PARALLEL GEOMETRIC MULTIGRID METHODS FOR POISSON EQUATION

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Abstract. Direct Numerical Simulation (DNS) of incompressible flows demands a high computational cost. The formulation adopted frequently requires the numerical solution of Poisson equations. Normally, this solution is the most intensive subroutine in the numerical code. Aiming the optimization of the solution of the Poisson equation, this paper deals with multigrid methods and parallelization. With this motivation in mind, and knowing that in theory, multigrid methods solve linear systems in computational time proportional to the dimensional grid, two parallel geometric multigrid methods, Correction Scheme (CS) and Full Approximation Scheme (FAS) for solving 2D Poisson equation were implemented and a performance analysis was conducted. The speedup results obtained indicate that, for coarse meshes, the values are significantly lower due to overhead. When the number of unknowns increases the speedup improves for both methods.

Keywords: Poisson equation, geometric multigrid, parallel algorithms.

1. INTRODUCTION

In recent years, significant advances in Numerical Analysis have been achieved, in order to optimize the time necessary to obtain solutions of fluid dynamic problems. These advances can be motivated, for example, by using Direct Numerical Simulation (DNS) of incompressible flows. Moreover, depending on the formulation adopted, the numerical solution of a Poisson equation is necessary. This solution is the most intensive subroutine since it depends on solving a linear system. Aiming the optimization of the solution of elliptic partial differential equations (PDE), a multigrid method can be used.

Historically, according to Wesseling (1992), multigrid method was first developed by Fedorenko in the 1960’s. Motivated by Fedorenko (1964), Brandt (1977) developed multigrid methods and studied the convergence rate, grid coarsening and local Fourier Analysis for linear and non-linear problems. Through Brandt’s contribution, the multigrid method became widely used by the scientific community to solve problems in fluid dynamics. Hirsch (1988) and Tannehill et al. (1997) claim that the multigrid method is considered one of the most efficient method proposed recently.

According to Van der Velde (1994), the benefits of using multigrid methods are justified because it combines iterative solvers in meshes with different number of discretization points. Furthermore, the works of Hirsch (1988) and Ferziger and Peric (1999) show that there is an independence between the number of iterations to achieve convergence in the finest mesh and the number of mesh points. Pavlov et al. (2001) and Wesseling (1984) multigrid methods solve an elliptic problem discretized with N points using N operations. Stüben (2001) presents advantages of the algebraic multigrid for unstructured meshes. Wesseling (2001) recommends the use of geometric multigrid for structured meshes. The present work aims to investigate geometric multigrid methods. It were implemented two parallel geometric multigrid methods, Correction Scheme (CS) and Full Approximation Scheme (FAS) for solving 2D Poisson equation and a performance analysis was conducted.

The paper is structured as follows: section 2 presents a description of the multigrid methods adopted in this work - CS and FAS - and their parallelization. In addition, this section describes the numerical approximations used. Section 3 presents a performance analysis of the implemented methods. Finally, the main conclusions, acknowledgments and references are described in sections 4, 5 and 6 respectively.

2. FORMULATION AND NUMERICAL METHODS

The mathematical equation adopted in the present work is the Poisson equation:

\[
\frac{\partial^2 v}{\partial x^2} + \frac{\partial^2 v}{\partial y^2} = s,
\]

where \( v \) represents the solution and \( s \) is the source term.

For the methods under consideration, a V-cycle composed of four uniform Cartesian grid meshes was used. Figure 1 represents this multigrid structure, where \( S \) represents the iterative/smooother method, \( R \) is the restriction operation and \( P \) is the prolongation operation. The finest grid is represented by \( h \) and the coarsest one by \( 8h \).

2.1 Correction Scheme

The idea of the CS method is associated with the estimation of the error generated in the finer mesh as a way to obtain the solution of the linear system. This estimation is obtained by transmitting the residual of a finer mesh to a coarser one.
The adopted method is illustrated in Fig. 2.

The CS multigrid method works as follow (Van der Velde, 1994):

1. Starting at the finest grid, two iterations of a Jacobi under-relaxed method with relaxation factor 0.9 is applied;

2. With the approximation of \( v \) the residual \( d_h \) can be achieved as:

\[
d_h = s_h - \nabla^2 v_h.
\]  \hspace{1cm} (2)

3. If the residual is smaller than a prefixed tolerance, the algorithm is ended. Otherwise, the residual \( d_h \) is transmitted to a coarse grid \((2h)\) through an operation called restriction. This operation is illustrated in Fig. 3.

\[
d_h \Rightarrow s_{2h} \quad (FW)
\]  \hspace{1cm} (3)

This specific operation is called Full Weight (FW). The only element to be transmitted, represented by a square, is defined as a source term in the coarse level \((2h)\) and it takes information of all neighbors using a specific weight for each one as shown in Fig. 3. The initial guess at \((2h)\) level is taken zero.

\[
v_{2h} \Rightarrow \text{corr}_{4h},
\]  \hspace{1cm} (4)

\[
v_{4h} \Leftarrow v_{4h} + \text{corr}_{4h}.
\]  \hspace{1cm} (5)

4. The return to a more refined mesh is done by an operation called prolongation using bilinear interpolation and then correcting the solution.
The bilinear interpolation is illustrated by Fig. 4. Each colored square represents a mesh point that exists only in the next refined grid. Moreover, the circles represent points existing in both levels. For blue and red squares, the information comes from the two inner circles and for the central green one, all four circles are used. The respective weights of the interpolation are shown above each arrow. This prolongation operation must be applied until the finest level. At this point, the cycle starts again.

\[ d_h = s_h - \nabla^2 v_h, \]  
\[ d_n \Rightarrow s_{2h} \quad (FW) \]  
\[ v_n \Rightarrow v_{2h} \quad (SI) \]  

The SI operation is illustrated in Fig. 6. Note that no information of the neighborhood is taken into account.

2.2 Full Approximation Scheme

In the present case the V-cycle structure can be illustrated by Fig. 5. The FAS multigrid method, described in Souza (2003), works as follow:

1. Starting at the finest grid, two iterations of SOR method with relaxation factor 1.1 are applied.
2. With this approximation of \( v \), the residual \( (d_h) \) is computed as:

\[ d_h = s_h - \nabla^2 v_h, \]  

3. If the residual is smaller than a tolerance, the algorithm is ended. Otherwise, the residual \( d_h \) is transmitted to a coarse grid \( (2h) \) through FW operation. Moreover, the approximate solution \( v_h \) is transmitted directly, without any kind of weighting, using Straight Injection operation (SI).

\[ d_h \Rightarrow s_{2h} \quad (FW) \]  
\[ v_n \Rightarrow v_{2h} \quad (SI) \]  

4. The correction is calculated by:

\[ corr_{2h} = v_{2h} - v_{2h}', \]  

where \( v_{2h}' \) represents the solution generated by the restriction operation and \( v_{2h}^n \) is the newest solution.
5. To return to a more refined level, first it is adopted the same prolongation operation used in the CS scheme and then the solution is corrected.

\[ \text{corr}_{sh} \Rightarrow \text{corr}_{4h}, \quad (10) \]
\[ v_{4h} \leftarrow v_{4h} + \text{corr}_{4h}. \quad (11) \]

6. Finally, one iteration of SOR method with factor 1.0 is applied. The V-cycle is finished after the iteration of SOR method at the finest grid.

2.3 Numerical approximations

The second-order derivatives in the Poisson equation were discretized using the following approximations (Lele, 1992; Souza, 2003). The \( i \) index means the grid position in the \( x \) direction, which ranges from 1 to \( imax \). In the same sense, \( j \) means the grid position in the \( y \) direction and ranges from 1 to \( jmax \).

\( x \) direction:
- for \( 2 < i < imax - 1 \):
  \[
  \frac{\partial^2 f}{\partial x^2} \bigg|_{i,j} = \frac{-f_{i-2,j} + 16f_{i-1,j} - 30f_{i,j} + 16f_{i+1,j} - f_{i+2,j}}{12\Delta x^2} + O(\Delta x^4). \quad (12)
  \]
- for \( i = 2 \):
  \[
  \frac{\partial^2 f}{\partial x^2} \bigg|_{2,j} = \frac{10f_{1,j} - 15f_{2,j} - 4f_{3,j} + 14f_{4,j} - 6f_{5,j} + f_{6,j}}{12\Delta x^2} + O(\Delta x^4). \quad (13)
  \]
- for \( i = imax - 1 \): a similar approximation of \( i = 2 \) was adopted.

\( y \) direction:
- for \( 2 < j < jmax - 1 \):
  \[
  \frac{2}{15} \frac{\partial^2 f}{\partial y^2} \bigg|_{i,j-1} + \frac{11}{15} \frac{\partial^2 f}{\partial y^2} \bigg|_{i,j} + \frac{2}{15} \frac{\partial^2 f}{\partial y^2} \bigg|_{i,j+1} =
  \frac{3f_{i,j-2} + 48f_{i,j-1} - 102f_{i,j} + 48f_{i,j+1} + 3f_{i,j+2}}{60\Delta y^2} + O(\Delta y^6). \quad (14)
  \]
- for \( j = 2 \):
  \[
  \frac{4}{5} \frac{\partial^2 f}{\partial y^2} \bigg|_{i,2} + \frac{1}{5} \frac{\partial^2 f}{\partial x^2} \bigg|_{i,3} = \frac{254f_{i,1} - 432f_{i,2} + 162f_{i,3} + 16f_{i,4}}{180\Delta y^2} + O(\Delta y^6). \quad (15)
  \]
- for \( j = jmax - 1 \): a similar approximation of \( j = 2 \) was adopted.

In the \( y \) direction, in problems with Neumann boundary conditions:
- for \( j = jmax - 1 \):
  \[
  \frac{\partial^2 f}{\partial y^2} = \frac{f_{i,jmax - 2} - 2f_{i,jmax - 1} + f_{i,jmax}}{\Delta y^2} + O(\Delta y^2). \quad (16)
  \]
• for \( j = j_{\text{max}} \):

\[
\frac{\partial^2 f}{\partial y^2} = -\frac{f_{i,j_{\text{max}}-2} + 8f_{i,j_{\text{max}}-1} - 7f_{i,j_{\text{max}}} - 2f_{i,j_{\text{max}}-1} + f_{i,j_{\text{max}}-2}}{2\Delta y^2} + 3\frac{\partial f}{\partial y} + O(\Delta y^2).
\]  

(17)

There is no need to adopt higher-order approximations near to \( j = j_{\text{max}} \), since the higher gradients are concentrated in the neighborhood of \( j = 1 \) in the fluid flow simulations that we focus in.

2.4 Parallelization of multigrid methods

The multigrid algorithms have been parallelized with Message Passing Interface (MPI) using a domain decomposition technique in the \( x \) direction. This choice is justified since the number of points in this direction is bigger than the number of points in the \( y \) direction, in the problems under consideration.

In this sense, considering a sequential rectangular mesh domain with \( imax \times j_{\text{max}} \) points. Each process is responsible for \( N_x \) points in the \( x \) direction and \( j_{\text{max}} \) points in the \( y \) direction. The value of \( N_x \) can be calculated by:

\[
N_x = \frac{imax + (\text{inter} + 1)(p - 1)}{p},
\]

(18)

\[
\text{inter} = 2^{(n-1)}(m - 1),
\]

(19)

where \( p \) represents the number of processes, \( \text{inter} \) is the superposition between adjacent subdomains, \( m \) is the length of the computational molecule adopted in the \( x \) direction and \( n \) is the number of a V-cycle’s levels.

For example, if a multigrid method using a V-cycle composed of 4 levels takes the following second-order approximation

\[
\frac{\partial^2 f}{\partial x^2} = \frac{f_{i+1,j} - 2f_{i,j} + f_{i-1,j}}{\Delta x^2} + O(\Delta x^2),
\]

(20)

then \( m = 3 \) and \( n = 4 \), which implies \( \text{inter} = 8 \) cells.

With regard to the algorithm’s parallelization, the same sequential multigrid methods were applied at each subdomain. However, since for elliptic problems there is a high dependence between the variables involved, it was necessary to introduce communication points between adjacent processes. For the CS method, these points are:

• before applying a restriction operation;
• at each step of the iterative/smooother method;
• before applying the interpolation.

For the FAS method, the communication occurs:

• after applying a restriction operation;
• at each step of the iterative/smooother method;
• before applying the interpolation.

The squares in Fig. 7 represent the communication points for both methods.

The algorithm for the parallelization of the CS scheme was found in the literature, but it was not found any description of parallel FAS algorithm, therefore the proposed parallel algorithm presents a contribution in this area.
3. NUMERICAL RESULTS

Tests were carried out with the Poisson equation using two different boundary conditions:

- case 1: Dirichlet condition in all contours.
- case 2: Neumann condition in one boundary and Dirichlet conditions in the others.

The evaluation was done considering 36 different Cartesian meshes using 1, 2, 4 and 8 processors. The processors used are Intel Xeon E5345 with 2.33GHz, 16GB RAM and 8 processing core. In order to verify the algorithms it was done an analysis of execution time and speedup. Speedup can be defined by:

\[ S_u = \frac{T_s}{T_p}, \]

where \( T_s \) represents the execution time for the algorithm with the lowest number of processes - for this case, a sequential program - and \( T_p \) represents the execution time considering the parallel program with \( p \) processes. The tolerance adopted in this work was \( 10^{-5} \). This criteria is associated with the absolute maximum residual of each subdomain.

The results obtained for execution time are shown in Fig. 8, where in the \( x \) direction, each point corresponds to a mesh in Tab. 1.

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Figure 8. Execution time

Analyzing the execution time, presented in Fig. 8, it is possible to see that when the number of variables increases, the CS method becomes better than FAS method, for both cases analyzed.

Furthermore, speedup for both cases with CS and FAS methods are shown in Figs. 9 and 10 respectively.

The speedup results obtained indicate that, for coarse meshes, the values are significantly lower due to overhead. In this sense, for meshes with small number of variables, sequential multigrid methods are better than parallel multigrid
methods. For example, considering the application of FAS/PSOR method with two processes, illustrated in Fig. 9(b), the use of sequential methods is better for the meshes 0 to 3 and 6 to 19. When the number of unknowns raises the speedup improves for both multigrid methods. The growth of speedup has an upper bound which is different for each case.

Considering the CS method, in both cases, the speedup’s growth behavior is similar for all values of $y$, considering the curve formed by fixing the number of points in $y$ direction and ranging the number of points in $x$ direction. This fact also can be seen in FAS/PSOR method until 257 points in $y$ direction. After that, the speedup rises rapidly with the increase in the number of points in $x$ direction.

For the first case, when the number of variables is small (for meshes from 0 to 25, for example), the CS method has a better speedup. While the number of unknowns increases, the speedup of FAS/PSOR method becomes similar to those of CS method. In addition, for the second case, in spite of being worse for small number of variables as in the first case, FAS/PSOR method achieves a better speedup than CS method, when the number of unknowns raises.

In general, considering just the first case, CS method showed the best results. Significant differences are observed for intermediate meshes. Nevertheless, for the second case, FAS method can be regarded as the best choice for finer meshes, as in cases 26-29 and 32-35.

4. CONCLUSIONS

Summarizing, for a large number of variables involved, the use of parallelism implies significant gains. For some cases this increase reaches more than five times over the sequential algorithm.

For the algorithms considered, the type of boundary condition used interferes in finding the best method. For the Dirichlet problem investigated, in general, CS method shows the best results. In the other hand, the use of a boundary condition of Neumann’s type implies the use of FAS/PSOR multigrid method for finer meshes.
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6. REFERENCES


7. Responsibility notice

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